

The Meta-ACG Preprocessor

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- 1 Introduction
- 2 Reducing redundancy
- 3 Simulating feature structures

My situation

My PhD goal : **ACG grammar** of French interrogatives in discourse and dialogue

- Using ACGtk

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- Using ACGtk
- Using **feature structures** (FS) :
 - agreement
 - anaphora resolution

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My PhD goal : **ACG grammar** of French interrogatives in discourse and dialogue

- Using ACGtk
- Using **feature structures** (FS) :
 - agreement
 - anaphora resolution
- Deep syntax inspired by HPSG [?]
 - SLASH feature for wh-phrase extraction

Why a preprocessor ?

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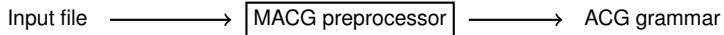
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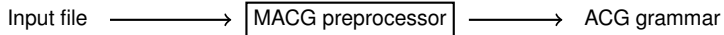
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Proposal : **Meta-ACG preprocessor** to automate this process

Building a preprocessor



Building a preprocessor



Usefulness :

- 1 Reduce redundancy
- 2 Simulate feature structures

ACGtk bureaucracy

Standard ACGtk grammar :

deep syntax

```
signature Deep =  
  n : type ;  
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(and similar for semantics)

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- Expected parts (conventional)
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- What really matters

Syntax of input file

MACG syntax :

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Command: param
        declaration
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```

Available commands :

- **Type** : mandatory param (the type name)
- **Constant**
- **Rule** : CFG rule

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Example :

```
Type: np
Type: vp
Type: s
```

Constant:

```
John, Mary : np
slept : vp
```

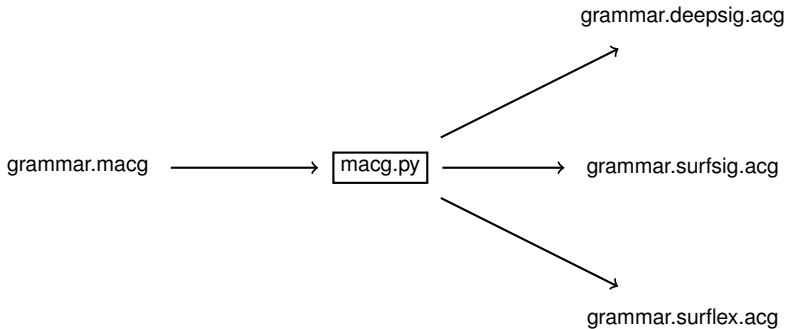
Rule:

```
np -> vp -> s
```

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- **Type** : mandatory param (the type name)
- **Constant**
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Preprocessing



Python program

Limitations

How much should be able to be customized ?

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 `SOME := lambda p, q. (Ex x. (p x) & (q x)) ;`
- ? Customizable signature / lexicon names

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 - Theoretical background
 - Solution

Typed feature structures

Feature structure (FS) = set of pairs **attribute-value**

$$\text{sleeps : } \left[\begin{array}{l} \textit{verb} \\ \text{AGR} \\ \text{TENSE} \end{array} \left[\begin{array}{l} \textit{agreement} \\ \text{PERSON} \quad 3 \\ \text{NUMBER} \quad \textit{singular} \\ \textit{present} \end{array} \right] \right] \quad (1)$$

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- **Typed** : each FS has predefined set of possible attributes and values

FS class :

$$\left[\begin{array}{l} \text{agreement} \\ \text{PERSON} \\ \text{NUMBER} \end{array} \left[\begin{array}{l} \{1, 2, 3\} \\ \{\textit{singular}, \textit{plural}\} \end{array} \right] \right] \quad (2)$$

FS instance :

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- **Underspecification** : some attributes may lack

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Subsumption and unification

Operations on FSs f, f'

f **subsumes** f' if they have the same type and f' specifies f more

e.g.

$$\left[\begin{array}{l} \textit{agreement} \\ \text{PERSON} \end{array} \quad 1 \right] \sqsubseteq \left[\begin{array}{l} \textit{agreement} \\ \text{PERSON} \\ \text{NUMBER} \end{array} \quad \begin{array}{l} 1 \\ \textit{singular} \end{array} \right],$$

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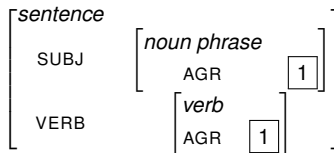
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but $\left[\begin{array}{l} \textit{agreement} \\ \text{PERSON} \end{array} \quad 1 \right]$ and $\left[\begin{array}{l} \textit{agreement} \\ \text{PERSON} \end{array} \quad 2 \right]$ can't unify

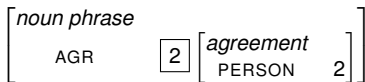
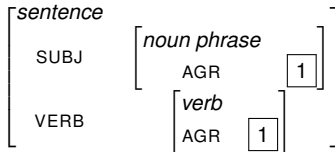
Unification in rules

E.g. The AGR values of the subject and the verb have to unify

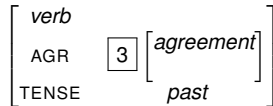


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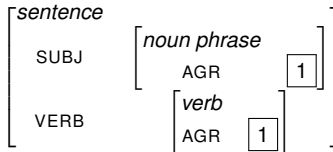
you



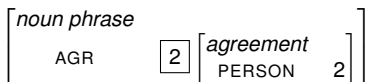
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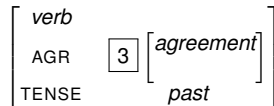
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Parsing is
finding $\boxed{1}$
 such that
 $\boxed{1} = \boxed{2} \sqcup \boxed{3}$



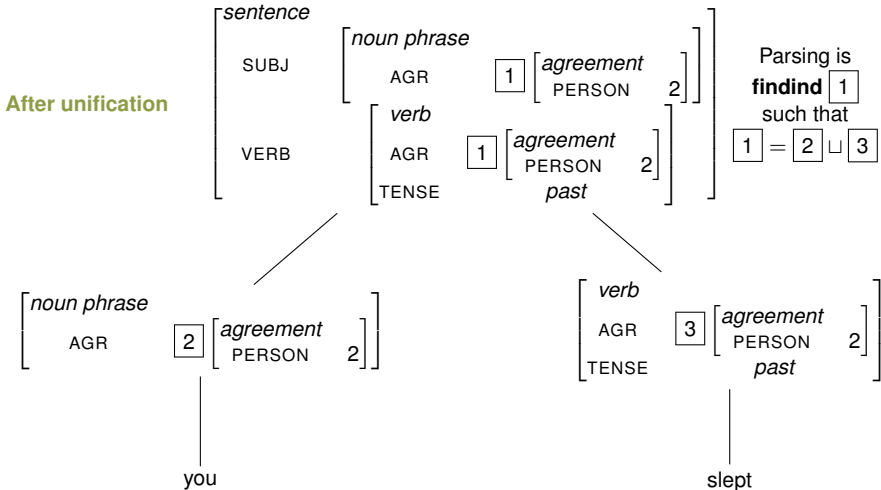
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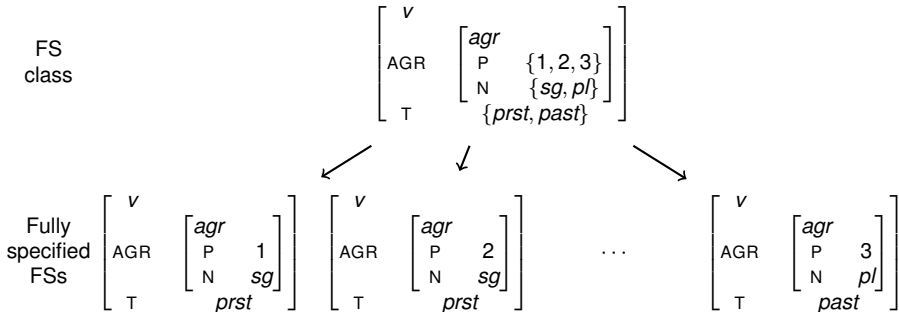
As many atomic types as possible FSs : **full instances**

FS
class

$$\left[\begin{array}{c} V \\ \text{AGR} \\ T \end{array} \left[\begin{array}{c} agr \\ P \quad \{1, 2, 3\} \\ N \quad \{sg, pl\} \\ \{prst, past\} \end{array} \right] \right]$$

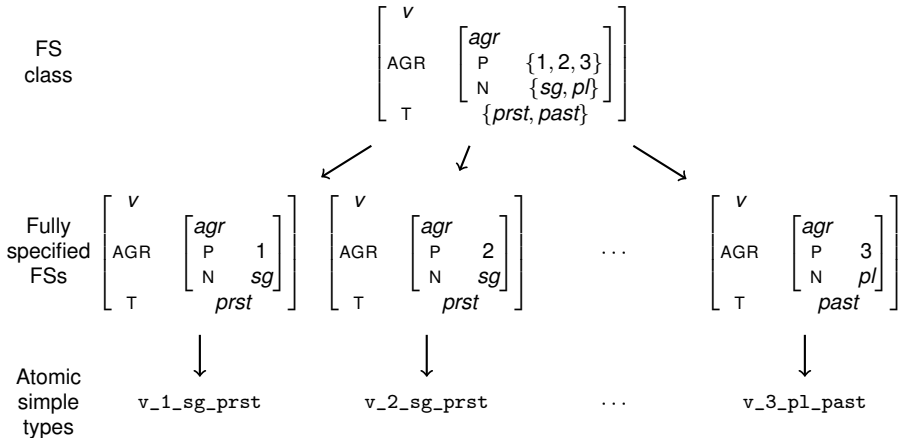
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MACG syntax like NLTK (python)

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Type classes must be defined before !

```
Type: agr
  P : 1, 2, 3
  N : sg, pl
Type: vp
  AGR : agr
  T : prst, past
Type: s
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Additional functionalities :

- boolean features, e.g. `is : v[+AUX]`
- Absence of a feature, e.g. no agreement `be : v[-AGR, T=inf]`

Formal proof of work

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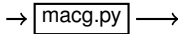
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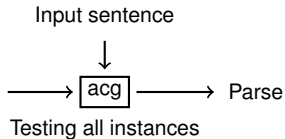
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Constants + rules
with FSs



All possible
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Future prospects :

- Higher-order types
- (Concatenative) morphological rules
- ? Type hierarchy and inheritance
- ? Reentrancy